## Scheme: Data

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## Topics

- ✓ PL feature: identifiers & values
- PL feature: reflection
- PL category: Lisp-family PLs
- Introduction to Scheme
- ✓ Scheme: basics
- ✓ Scheme: evaluation
  - Scheme: data
  - Scheme: macros

## Review

Scheme is a Lisp-family PL with a minimalist design philosophy.

Scheme code consists of parenthesized lists, which may contain atoms or other lists. List items are separated by space; blanks and newlines between list items are treated the same.

```
(define (print-sum-2-7)
  (display (+ 2 7))
)
```

Normal evaluation rule for a list: attempt to evaluate each list item, then attempt to call the result from the first item, as a **procedure**, with the results from the others as its arguments. For example, display and + (above) evaluate to procedures. Things that break this rule, like define (above), are **macros**. *Numeric* comparison operators: = < <= > >=

There is no standard numeric inequality operator.

Use these comparison operators only with numbers.

Given a nonempty list, car returns its first item (head), and cdr returns a list of the remaining items (tail). cons constructs a list, given head and tail.

```
> (car '(1 2 3))
1
> (cdr '(1 2 3))
(2 3)
> (cons 1 '(2 3))
(1 2 3)
```

Actually, car, cdr, and cons are more general, working with pairs, not just lists. More on this soon. eval is a procedure that takes one argument and evaluates it.
Being a procedure, eval does not suppress normal argument evaluation. So evaluation actually happens twice: the argument is evaluated, and then it evaluates the result.



A variation is the procedure apply. This takes a procedure and a list of arguments. It calls the procedure with the given arguments and returns the result.

```
> (apply + '(2 3))
```

5

## Scheme: Data

The dot notation originally used in S-expressions is also valid in Scheme.

> '(1 . 2) (1 . 2) For code from this topic, see data.scm.

List notation is really shorthand for the equivalent dot notation, again, just as in the original S-expression syntax.

```
> '(1 . (2 . (3 . (4 . ()))))
(1 2 3 4)
```

Dot (.) is not a procedure. It is simply another way of typing Sexpressions. If you want a procedure that puts things together the way dot does, use cons. Scheme: Data Data Format [2/5]

The main building block for constructing data structures in Scheme is the **pair**. You can think of this as a node with two pointers.

$$(1 \cdot 2) \begin{cases} 1 \\ 1 \\ 1 \\ 2 \end{cases}$$

We get the item referenced by the left pointer using **car**; similarly use **cdr** for the right pointer.

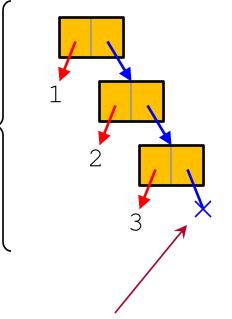
```
> (car '(1 . 2))
1
> (cdr '(1 . 2))
2
```

Scheme: Data Data Format [3/5]

Lists are constructed from pairs and null.

$$(1 \ 2 \ 3) \equiv (1 \ . \ (2 \ . \ (3 \ . \ ())))$$

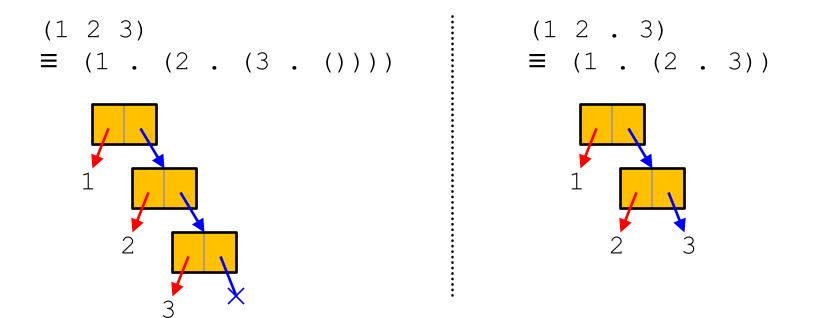
I am using this symbol to mean "is the same as".



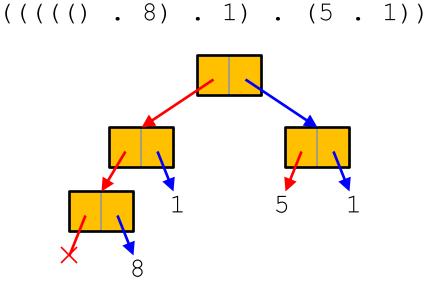
This represents *null*. Think of a null pointer, if you want. (But how it is represented internally is implementation-dependent.)

Scheme: Data Data Format [4/5]

The full story on the dot syntax is that the dot may optionally be added between the *last two* items of something that otherwise looks like a list.



We can create arbitrary binary trees—with the restriction that only leaf nodes contain data.



Scheme has type-specific comparison procedures.

Comparisons for numbers, as we have seen: = < <= > >=

```
> (= 3 3.0)
#t
> (> 3 3.1)
#f
```

Comparisons for non-numeric types are named as type+op+?

```
> (string=? "abc" "abc")
#t
> (char<? #\b #\a)
#f</pre>
```

There are several kinds of equality in Scheme. eq? tests for "same memory location". I suggest *not* using it. eqv? tests for "same primitive value".

```
> (eqv? 3 3)
#t
> (eqv? 3 3.0)
#f
> (define a '(1 2))
> (eqv? a '(1 2))
#f
> (eqv? a '(1 2))
Lists and strings are
not primitive values.
> (eqv? "abc" "abc")
IMPLEMENTATION-DEPENDENT
```

It is common to use eqv? indirectly. See the next slide.2025-03-26CS 331 Spring 2025

Of greater interest is equal?, which does the following:

- If the types are different, then return #f.
- For primitive values (everything we have covered except strings and pairs) of the same type, call eqv?.
- For pairs, recursively call equal? on the cars & cdrs.
- For other non-primitive values (e.g., strings) of the same type, call an appropriate type-specific equality comparison, if one exists.

So for lists, equal? checks structural equality.

```
> (define a '(1 (2 3) 4))
> (equal? a '(1 (2 3) 4))
#t
> (equal? "abc" "abc")
#t
```

equal? mostly does what we usually want, with one caveat. Since it always returns #f when the types are different, it can give undesirable results with numbers.

```
> (equal? 3 3.0)
#f
```

I offer the following suggestions.

- Use = for numeric equality.
- Using equal? is fine for most other kinds of equality.
- If you want your code to indicate what type is being compared, or to flag type errors for other types, then use a type-specific equality function: string=?, char=?, etc.
- Use eq? or eqv? directly only if you are sure of what you are doing—and probably never.

So far, all the procedures we have written have taken a fixed number of parameters. But Scheme allows for procedures like "+", which can take an arbitrary number of parameters.

Let's duplicate "+", in the form of a procedure called add.

```
> (add 5 3)
8
> (add 1 2 3 4)
10
> (add)
0
```

We will use "+", but only as a 2-parameter procedure.

Consider a call to add:

```
> (add 1 2 3 4)
10
```

The above list (add 1 2 3 4) is the same as (add . (1 2 3 4)).

So a procedure call is a pair. The car is the procedure; the cdr is a *list* of the arguments. This is illustrated below.



Scheme: Data More General Procedures [3/4]

A procedure call is a pair: (*PROC* . *ARGS*). And define will also take this form of a "picture" of a procedure call.

```
(define (add . args)
...
)
TO DO
```

• Write procedure add.

Done. See data.scm.

A tricky issue is how to make a recursive call on (cdr args). We look at this on the next slide.

In writing procedure add, we need to make a recursive call on (cdr args). How do we do this?

NOT like this (dot is not a procedure!):

(add . (cdr args))
is just another way to write
(add cdr args), which is
not what we want.

The following will actually work, but it is a bit unwieldy:

; WRONG!

```
(eval (cons add (cdr args)))
```

(edr args))

Situations like this are why apply exists:

(apply add (cdr args))

(add

We can write code that deals with a structure, not as a list, but as a tree, traversing the tree and dealing with atoms in some way.

TO DO

- Write a procedure atomsum that is given a tree t and returns the sum of all the numbers in t.
- Write a procedure atommap that is given a procedure f and a tree t and returns t with each atom replaced by f of that atom.
- Write a procedure myflatten that is given a tree t and returns a list of the atoms in t in inorder-traversal order.
- In order to write myflatten easily, we need to be able to concatenate two lists. Write this first, as procedure concat.

Done. See data.scm.

May be helpful:

Every Scheme value is null or a pair or an atom. So any value for which both null? and pair? both return #f is an atom.

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CS 331 Spring 2025