Thoughts on Project 7
Binary Heap Algorithms
Heaps & Priority Queues in the C++ STL

CS 311 Data Structures and Algorithms
Lecture Slides
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Thoughts on Project 7
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Overview

In Project 7, you will implement the *Treesort* algorithm.

```cpp
template<typename FDIter>
void treesort(FDIter first, FDIter last);
```

Treesort is a general-purpose comparison sort that uses a Binary Search Tree—which you will need to implement.

We look at:

- How Treesort works
- Things you will need to do
  - Writing a Binary Search Tree
  - Finding the value type of an iterator
  - Inserting into a Binary Search Tree
  - Inorder traversal of a Binary Search Tree
Thoughts on Project 7
Treesort — Introduction

A sorted container often gives us a sorting algorithm: insert all items into the container, and then iterate through it.

For a SortedSequence, this algorithm is *almost* Insertion Sort. (It would be a non-in-place version of Insertion Sort.)

For a Binary Search Tree, the algorithm is called **Treesort**.

- **Procedure**
  - Go through the list to be sorted, inserting each item into a BST.
  - Traverse (inorder) the tree; copy each item back to the original list.
- We must allow multiple equivalent keys in our BST.
- Treesort is not a terribly good algorithm.
Treesort this list.

- Go through the list to be sorted, inserting each item into a BST.

- Traverse (inorder) the tree; copy each item back to the original list.
Thoughts on Project 7
Treesort — Properties

What is the order of Treesort?

- Treesort does \( n \) BST inserts; each is \( \Theta(n) \). Then a traverse: \( \Theta(n) \).
- So Treesort is \( \Theta(n^2) \).
- However, BST insert is \( \Theta(\log n) \) on average.
- So Treesort is pretty fast on average: \( \Theta(n \log n) \).

Q. Have we seen Treesort before?
A. Kind of. It is much like an unoptimized Quicksort—in disguise.

Two important differences:

- Treesort requires a large additional memory space, because it does not do its work in the original storage.
- Treesort is not limited by the requirements of a fast in-place partition algorithm. As a result, it is stable.
Thoughts on Project 7
Writing a Binary Search Tree [1/2]

You will need to define a Binary Search Tree node type (class/struct—or use `std::tuple`).

It is *not* necessary to write a tree class.
- A tree can be handled via a (smart?) pointer to its root node.
- However, you may write a tree class if you wish.

In the following slides, I assume you use a simple struct for a Binary Search Tree node, and you refer to a node with a `unique_ptr`. But do things differently if you want.
Your Binary Search Tree implementation does not need to include all BST operations.

Treesort does the following:

1. Insert all items into a Binary Search Tree.
2. Traverse the tree, copying items to the original range.
3. Destroy the tree.

Destruction is automatic, if you use smart pointers. So you need to be able to insert an item and do an inorder traversal.
Thoughts on Project 7
Finding the Value Type

The type of the values to be sorted can be determined as follows.

```cpp
#include <iterator>
// For std::iterator_traits;

template<typename FDIter>
void treesort(FDIter first, FDIter last)
{
    using Value = typename
        std::iterator_traits<FDIter>::value_type;

    Then you can do things like this:

    auto p = std::make_unique<BSTreeNode<Value>>(...
```

This is my node struct.
Thoughts on Project 7
Binary Search Tree Insert [1/2]

*Suggestion.* Write a Binary Search Tree insert function that takes:
- A smart pointer to a tree node, *by reference*.
- The item to insert.

**Operation:**
- If the pointer is null/empty, then set it to point to a new node.
- Otherwise, it points to a node. Compare that node’s item with the given item. *Recurse* with the node’s left- or right-child pointer, as appropriate.

To insert an item into the tree, call the above function with:
- The head pointer for the tree.
- The item to be inserted.
So your insert function would look like the following.

```cpp
template<typename Value>
void insert(unique_ptr<BSTreeNode<Value>> & head, 
            const Value & item);
```
Thoughts on Project 7
Inorder Traversal

*Suggestion*. Write an inorder traversal function that takes:
- A pointer to a tree or node.
- An iterator, *by reference*.

**Operation:**
- If the pointer is null, return.
- Recurse with the left-child pointer.
- **Write the data to the location referenced by the iterator.**
- Increment the iterator.
- Recurse with the right-child pointer.  

To copy the data in the tree back to the original storage, call the above function with:
- The head pointer for the tree.
- Parameter *first* from function tree_sort.
Review
Review
Where Are We? — The Big Challenge

Our problem for most of the rest of the semester:

- Store: A collection of data items, all of the same type.
- Operations:
  - Access items [single item: retrieve/find, all items: traverse].
  - Add new item [insert].
  - Eliminate existing item [delete].
- Time & space efficiency are desirable.

A solution to this problem is a **container**.
In a **generic container**, client code can specify the value type.
Unit Overview
Tables & Priority Queues

Major Topics
- Introduction to Tables
- Priority Queues
- Binary Heap Algorithms
  - Heaps & Priority Queues in the C++ STL
  - 2-3 Trees
  - Other self-balancing search trees
- Hash Tables
- Prefix Trees
- Tables in the C++ STL & Elsewhere
A **Table** allows for look-up by arbitrary key.

Three primary operations: retrieve, insert, delete.

Possible Table implementations (all too slow):

- A Sequence holding key-value pairs.
  - Array-based or Linked-List-based.
  - Sorted or unsorted.
- A Binary Search Tree holding key-value pairs.
  - Implemented using a pointer-based Binary Tree.

### Array Implementations

- **Unsorted**
  - (12, Ed)
  - (4, Peg)
  - (9, Ann)

- **Sorted**
  - (4, Peg)
  - (9, Ann)
  - (12, Ed)

### Linked List Implementations

- **Unsorted**
  - (12, Ed)
  - (4, Peg)
  - (9, Ann)

- **Sorted**
  - (4, Peg)
  - (9, Ann)
  - (12, Ed)

### Binary Search Tree Implementation

- (9, Ann)
- (4, Peg)
- (12, Ed)

### Table

<table>
<thead>
<tr>
<th>Key</th>
<th>Value</th>
</tr>
</thead>
<tbody>
<tr>
<td>12</td>
<td>Ed</td>
</tr>
<tr>
<td>4</td>
<td>Peg</td>
</tr>
<tr>
<td>9</td>
<td>Ann</td>
</tr>
</tbody>
</table>
Idea #1: Restricted Tables
- Only allow retrieve/delete on the *greatest* key.
- In practice: Priority Queues

Idea #2: Keep a tree balanced
- In practice: Self-balancing search trees (2-3 Trees, etc.)

Idea #3: Magic functions
- Use an unsorted array. Each item can be a key-value pair or *empty*.
- A *magic function* tells the index where a given key is stored.
- Retrieve/insert/delete in constant time? No, but still a useful idea.
- In practice: Hash Tables
Unit Overview
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- Binary Heap Algorithms
  - Heaps & Priority Queues in the C++ STL
  - 2-3 Trees
  - Other self-balancing search trees
- Hash Tables
- Prefix Trees
- Tables in the C++ STL & Elsewhere

Lots of lousy implementations
Idea #1: Restricted Table
Idea #2: Keep a tree balanced
Idea #3: Magic functions
A special-purpose implementation: “the Radix Sort of Table implementations”
A **Binary Heap** (or just **Heap**) is a complete Binary Tree with one data item—which includes a **key**—in each node, where no node has a key that is less than the key in either of its children.

In practice, we use “Heap” to refer to the array-based complete Binary Tree implementation of this.
A Heap is a good basis for an implementation of a **Priority Queue**.

- Like Table, but retrieve/delete only highest key.
- Insert any key-value pair.

Algorithms for three primary operations

- **getFront**
  - Get the root node.
  - Constant time.

- **insert**
  - Add new node to end of Heap. Sift-up last item.
  - Logarithmic time if space available.
  - Linear time if not. However, in practice, a Heap often does not manage its own memory.

- **delete**
  - Swap first & last items. Reduce size of Heap. Sift-down new root item.
  - Logarithmic time. **Faster than linear time!**
Review
Binary Heap Algorithms [3/8]

To insert into a Heap, add new node/item at end. Then **sift-up**.

- If value is in root, or is $\leq$ its parent, then stop.
- Otherwise, swap item with parent. Repeat at new position.

```
56
  50
   5
    1 3 10 3
  22
   25 11
    12 32
```

```
56
  50
   5
    1 3 10 3
  22
   25 11
    12 25
```

```
56
  50
   5
    1 3 10 3
  32
   11
    12 25
```

```
56
  50
   5
    1 3 10 3
  22
   25 11
    12 25
```

Done

Ick!
To delete the root item from a Heap, swap root & last items, and reduce the size of the Heap. The **sift-down** the new root item.

- If value is ≥ all of its children, the stop.
- Otherwise, swap item with its largest child. Repeat at new position.
Do Heap insert 30 on the Binary Heap shown below. Draw the resulting Heap, as a tree.

Answer on next slide.
Review
Binary Heap Algorithms [6/8] (Try It!)

Do Heap insert 30 on the Binary Heap shown below. Draw the resulting Heap, as a tree.
Heap insert and delete are usually given a random-access range. The item to insert or delete is the last item; the rest is a Heap.

- **Action of Heap insert:**
  - This is where we want the item to be, initially (right?).

- **Action of Heap delete:**
  - This is where the swap puts the item (right?).

Note that Heap algorithms can do *all* modifications using *swap*. This usually allows for both speed and (exception) safety.
TO DO

- Write the Heap insert algorithm.
  - Prototype is shown below.
  - The item to be inserted is the final item in the given range.
  - All other items should form a Heap already.

```cpp
// Requirements on types:
//   RAIter is a random-access iterator type.

template<typename RAIter>
void heapInsert(RAIter first, RAIter last);
```

Done. See heap_algs.h. The other basic Heap algorithms have also been implemented.

See heap_algs_main.cpp for a program that uses this header.
Binary Heap Algorithms

continued
Binary Heap Algorithms
Fast Make-Heap

To turn a range (array?) into a Heap, we could do $n-1$ Heap inserts.
- Each insert operation is $\Theta(\log n)$; making a Heap in this way is $\Theta(n \log n)$.

However, there is a faster way.
- Place each item into a partially-made Heap, in \emph{backwards order}.
- Sift-down each item through its descendants.
  - For most items, there are not very many of these.

This Make-Heap algorithm is \emph{linear time}!
Our last sorting algorithm is **Heap Sort**.

- This is a fast comparison sort that uses Heap algorithms.
- We can think of it as using a Priority Queue, except that the algorithm is in-place, with no separate data structure used.
- Procedure. Make a Heap, then delete all items, using the Heap delete procedure that places the deleted item in the top spot.

**Is Heap Sort efficient?**

- The **Make-Heap** operation is $\Theta(n)$. Then we do $n$ **Heap delete** operations, each of which is $\Theta(\log n)$.
- Total: $\Theta(n \log n)$.

---

*See heap_sort.cpp for an implementation of Heap Sort. This uses the Heap algorithms in heap_algs.h.*
Heap Sort is in-place.

- We can create a Heap in a given array.
- As each item is removed from the Heap, put it in the array item that was removed from the Heap.
  - Starting the delete by swapping root and last items does this.
- Results
  - Ascending order, if we used a Maxheap.
  - Only constant additional memory is required. No reallocation is done.

So Heap Sort uses less space than Introsort or array Merge Sort.

- Heap Sort: $\Theta(1)$.
- Introsort: $\Theta(\log n)$.
- Merge Sort on an array: $\Theta(n)$.

Heap Sort can easily be generalized.

- Stopping before the sort is finished.
- Doing Heap insert operations in the middle of the sort.
Below: Make-Heap operation. Next slide: Heap deletion phase.

Start

Add 4

Add 1

This is what happens in memory.

This is a picture of the logical structure.

Now the entire array is a Heap.
Heap deletion phase:

Start

Delete 4

Delete 3

Delete 2

The last Heap delete does nothing, and may be omitted.

Now the array is sorted.
Binary Heap Algorithms
Heap Sort — Analysis

Efficiency 😊
- Heap Sort is \( \Theta(n \log n) \).

Requirements on Data 😞
- Heap Sort requires random-access data.

Space Usage 😊
- Heap Sort is in-place.

Stability 😞
- Heap Sort is not stable.

Performance on Nearly Sorted Data 😞
- Heap Sort is not significantly faster or slower for nearly sorted data.

Notes
- Heap Sort can be stopped early, with useful results.
- Recall that Heap Sort is used by Introsort, when the depth of the Quicksort recursion exceeds the maximum allowed.

We have seen these together before (Iterative Merge Sort on a Linked List), but never for an array.
In practice, a Binary Heap is not so much a data structure as it is a random-access range with a particular ordering property.

Associated with Heaps are a collection of algorithms that allow us to efficiently create Priority Queues and do comparison sorting.

- These algorithms are the things to remember.
- Thus the subject heading.
Heaps & Priority Queues in the C++ STL
The C++ STL includes several Heap algorithms, in `<algorithm>`.  
- Each takes a range specified by a pair of random-access iterators.  
- An optional third parameter is a custom comparison.

```cpp
std::push_heap
- Heap insert. `[first, last-1)` is a Heap. Item to insert is `*(last-1)`.
```
```
std::pop_heap
- Heap delete. Puts the deleted element in `*(last-1)`.
```
```
std::make_heap
- Make a range into a Heap, using fast Make-Heap algorithm.
```
```
std::sort_heap
```
```
std::is_heap
- Test whether a range is a Heap. Returns `bool`.
```

Calling `make_heap` and then `sort_heap` on the same range, will sort the range using Heap Sort.
std::partial_sort, in `<algorithm>`, does Heap Sort, but may stop early.

- It takes three iterators: `first`, `middle`, `last`, and an optional comparison.
- `[first, last)` must be a valid range, and `middle` must be between `first` & `last` (inclusive).
- It does Heap Sort—in reverse order using a Minheap—stopping when the range `[first, middle)` has been filled with sorted data.
- Result. The range `[first, middle)` contains exactly the data it would if the entire range were sorted. The range `[middle, last)` contains the items it would contain if the entire range were sorted, but they may not be in sorted order.

A variation, `std::partial_sort_copy`, also in `<algorithm>`, leaves data in the original range unchanged, placing its results in a second provided range.
The STL has a Priority Queue: `std::priority_queue (<queue>)`. This is another *container adapter*: wrapper around a container. And once again, you get to pick what that container is.

```cpp
std::priority_queue<T, container<T>>
```

- `T` is the value type.
- `container<T>` can be any standard-conforming *random-access* sequence container with value type `T`.
- In particular, `container` can be `vector` or `deque`.
- But not `list`, as it is not *random-access*.

`container` defaults to `std::vector`.

```cpp
std::priority_queue<T>  // = std::priority_queue<T, std::vector<T>>
```
Heaps & Priority Queues in the C++ STL
std::priority_queue — Members

The member function names used by std::priority_queue are the same as those used by std::stack.

- Not those used by std::queue.
- So std::priority_queue has top, not front.

Given a variable pq of type std::priority_queue<T>, we can do:

- pq.top()
- pq.push(item)
  - item is some value of type T.
- pq.pop()
- pq.empty()
- pq.size()
The comparison used by `std::priority_queue` defaults to `operator<`.

```cpp
std::priority_queue<Foo> pq1;  // Uses operator<
```

We can specify a custom comparison.
- An optional template parameter is the `type` of a comparison object.

Below is a Priority Queue whose `top` value is the smallest.

```cpp
std::priority_queue<Foo, std::vector<Foo>, std::greater<Foo>> pq2;
  // Uses operator>
```

We give the third template argument, so we must also give the second.

If, as above, we pass no ctor arguments, then our comparison object is a default-constructed object of the given type.
To pass our own comparison function, we specify:
- Its type, as the third template argument.
- The comparison itself, as a constructor argument.

```cpp
auto comp = [] (const Foo & a, const Foo & b) {
    return a.bar() < b.bar();
};
```

```cpp
std::priority_queue<Foo, std::vector<Foo>, decltype(comp)> pq3(comp);
```

See pq.cpp for example code.
We will look at a more practical example near the end of the semester.
Overview of Advanced Table Implementations

This ends our coverage of Idea #1: restricted Tables. Next, actual Tables, allowing retrieve & delete for arbitrary keys. We cover the following advanced Table implementations.

- Self-balancing search trees
  - To make things easier, allow more children (?):
    - **2-3 Tree**
      - Up to 3 children
    - **2-3-4 Tree**
      - Up to 4 children
    - **Red-Black Tree**
      - Binary Tree representation of a 2-3-4 Tree
    - Or back up and try for a strongly balanced Binary Search Tree again:
      - **AVL Tree**

- Alternatively, forget about trees entirely:
  - **Hash Table**

- Finally, “the Radix Sort of Table implementations”:
  - **Prefix Tree**

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Idea #2: Keep a tree balanced

Later, we will cover other self-balancing search trees: B-Trees, B+ Trees.

Idea #3: Magic functions