

Scheme: Strings & I/O

CS F331 Programming Languages

CSCE A331 Programming Language Concepts

Lecture Slides

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Scheme is a Lisp-family PL with a minimalist design philosophy.

Scheme code consists of parenthesized lists, which may contain atoms or other lists. List items are separated by space; blanks and newlines between list items are treated the same.

```
(define (hello-world)
  (begin
    (display "Hello, world!")
    (newline)
  )
)
```

When a list is evaluated, the first item should be a **procedure** (think “function”); the remaining items are its arguments.

The type system of Scheme is similar to that of Lua.

- Typing is dynamic.
- Typing is implicit. Type annotations are generally not used.
- Type checking is structural. Duck typing is used.
- There is a high level of type safety: operations on invalid types are not allowed, and implicit type conversions are rare.
- There is a fixed set of types (36 of them).

Scheme has 36 types (as compared to Lua's 8).

`quote` is a special procedure that takes one parameter, suppressing the normal parameter evaluation. It returns this parameter.

```
> (quote (1 2 3))  
(1 2 3)
```

The leading-single-quote syntax is actually shorthand for `quote`.

```
> '(1 2 3) ; Same as (quote (1 2 3))  
(1 2 3)
```

*For code from this topic,
see `data.scm`.*

`eval` is a procedure that takes one parameter and evaluates it.

`eval` does not suppress the normal evaluation of parameters, so, strictly speaking, evaluation happens twice: the parameter is evaluated, and then it evaluates the result.

```
> (eval '(+ 2 3))
```

5

A variation is `apply`. This takes a procedure and a list of arguments. It calls the procedure with the given arguments and returns the result.

```
> (apply + '(2 3))
```

5

The dot notation originally used in S-expressions is also valid in Scheme.

```
> '(1 . 2)  
(1 . 2)
```

A list is really shorthand for the equivalent dot notation, again, just as in the original S-expression syntax.

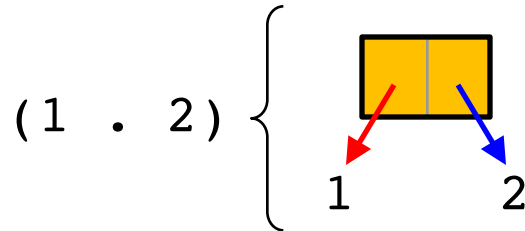
```
> '(1 . (2 . (3 . (4 . ())))))  
(1 2 3 4)
```

Dot (.) is *not a procedure*. It is simply another way of typing S-expressions. If you want a procedure that puts things together the way dot does, use `cons`.

Review

Scheme: Data — Data Format [2/5]

The main building block for constructing data structures in Scheme is the **pair**. You can think of this as a node with two pointers.



We get the item referenced by the left pointer using `car`; similarly use `cdr` for the right pointer.

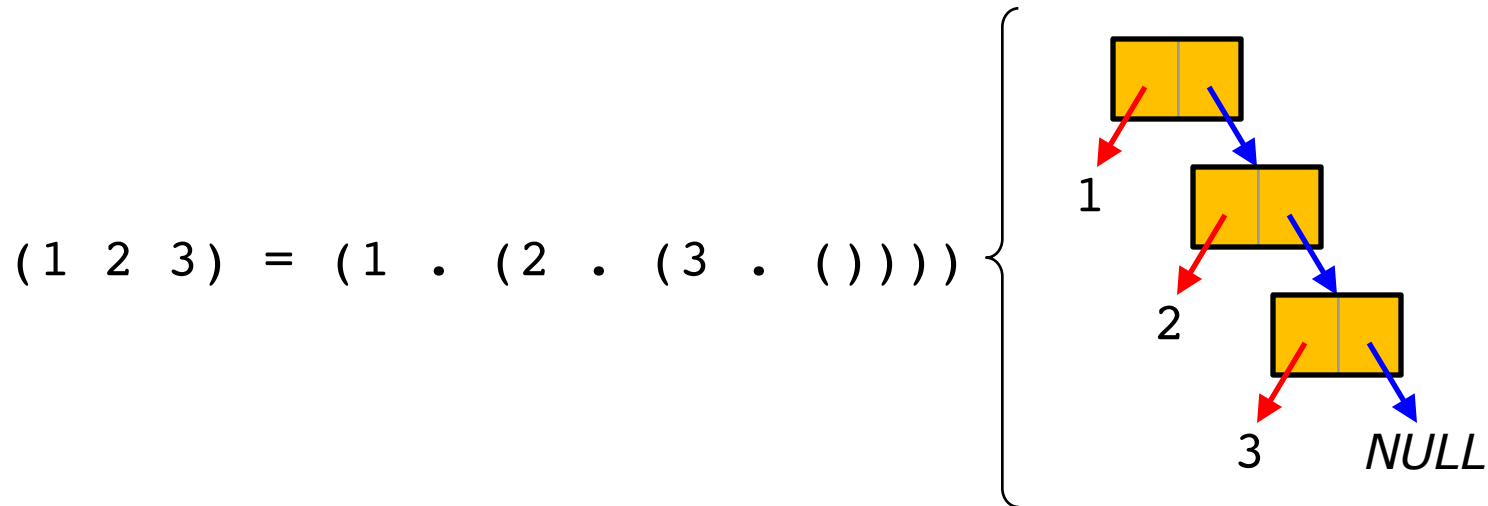
```
> (car '(1 . 2))
```

```
1
```

```
> (cdr '(1 . 2))
```

```
2
```

Lists are constructed from pairs and null.

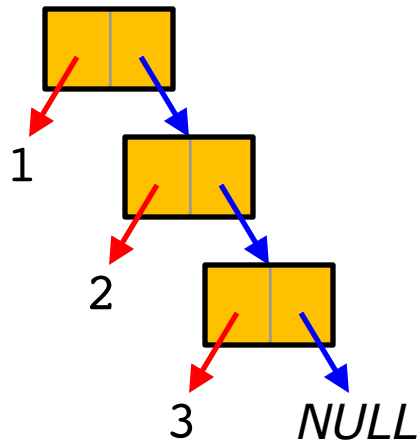


Review

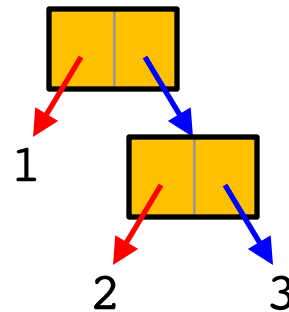
Scheme: Data — Data Format [4/5]

The full story on the dot syntax is that the dot may optionally be added just before the *last* item of a list.

(1 2 3)
= (1 . (2 . (3 . ())))



(1 2 . 3)
= (1 . (2 . 3))

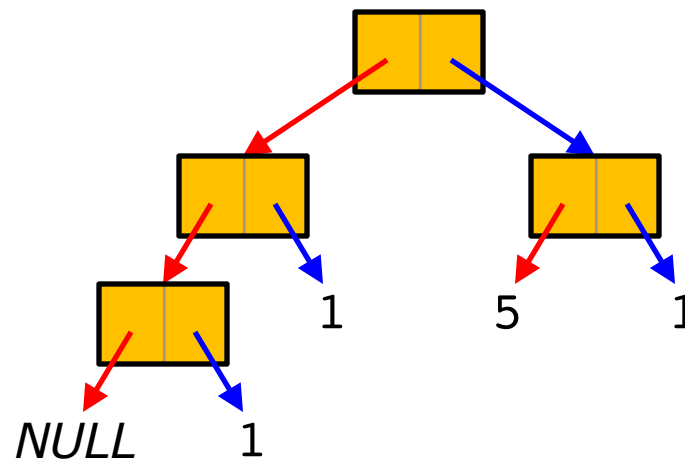


Review

Scheme: Data — Data Format [5/5]

We can create arbitrary binary trees—with the restriction that only leaf nodes contain data.

((((() . 8) . 1) . (5 . 1))



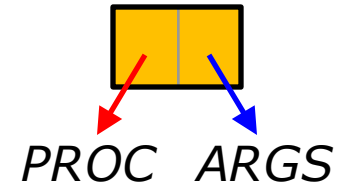
Review

Scheme: Data — Varying Number of Parameters

A procedure call is a pair: (*PROC* . *ARGS*).

`define` will take this form of a “picture” of a procedure call.

Procedure Call



```
(define (sum . args)
  ...
)
```

← *args* is a list of the arguments of *sum*. So *sum* can take an arbitrary number of parameters.

How to make a recursive call on (*cdr* *args*)?

~~(sum . (cdr args))~~

; WRONG!

(eval (cons sum (cdr args)))

; Okay

(apply sum (cdr args))

; Okay (and also clearer)

← (sum . (cdr args)) is just another way to write (sum cdr args), which is not what we want.

Normal evaluation in Scheme is eager.

However, we can do lazy evaluation in Scheme, using a **promise**: a wrapper around an expression that leaves the expression unevaluated.

When we **force** a promise, the expression is evaluated, and the resulting value is stored in the promise and returned. Force again, and the same value is returned, without reevaluating the expression.

Create a promise using `delay`.

```
> (define pp (delay (* 20 5)))
```

The type-checking predicate for promises is `promise?`.

```
> (promise? pp)
```

```
#t
```

Force a promise using `force`. Again, force a promise as many times as you like; evaluation only happens the first time.

```
> (force pp)
```

```
100
```

Using promises, we can create the kind of lazy infinite lists we saw in Haskell (rather less conveniently, though).

One way to do this is to construct a list as usual, from pairs and null, but wherever there is a pair, we actually have a promise wrapping a pair.

TO DO

- Write code to create a lazy infinite list.
- Write code to print out a portion of the above list. (With a little thought, we can write a procedure that will print both lazy lists and normal lists.)

See data.scm.

Scheme: Strings & I/O

Strings [1/3]

As in so many PLs, to understand Scheme I/O, it helps to know something about Scheme strings.

String literals in Scheme are surrounded by double quotes. The usual backslash escapes are accepted.

```
"This is a string."
```

```
"A newline: \nA double quote: \" A backslash: \\"
```

Check whether a value is a string with `string?`.

```
> (string? "42")
```

```
#t
```

```
> (string? 42)
```

```
#f
```

*For code from this topic,
see `string.scm`.*

Scheme: Strings & I/O

Strings [2/3]

Get the length of a string with `string-length`.

```
> (string-length "Hello!")  
6
```

Concatenate strings with `string-append`.

```
> (string-append "abc" "def" "ghi" "jklmnop")  
"abcdefghijklmnop"
```

Get a substring with `(substring STRING START PAST_END)`.

```
> (substring "Howdy thar!" 2 7) ; Zero-based indices  
"wdy t" ← Includes the characters at indices 2, 3, 4, 5, 6, but not 7.
```


Scheme: Strings & I/O

Strings [3/3]

Convert a number to a string using `number->string`.

```
> (number->string 42)  
"42"
```

Convert a string to a number using `string->number`. This returns the number, or `#f` if the conversion could not be done. So the result can be used in an `if`.

```
> (string->number "42")  
42  
> (string->number "Hello!")  
#f
```

(if COND THEN-EXPR ELSE-EXPR)

When the above is evaluated,
THEN-EXPR is chosen if *COND*
evaluates to *anything* other than *#f*.

Scheme: Strings & I/O

Characters [1/2]

Character literals generally have the form `#\CHAR`. Some characters have special literals.

`#\A` ; The 'A' character

`#\newline` `#\space`

Check whether a value is a character with `char?`.

```
> (char? #\x)
```

```
#t
```

```
> (char? "x")
```

```
#f
```

```
> (string? #\x) ; A Scheme character is not a string
```

```
#f
```

Scheme: Strings & I/O

Characters [2/2]

Convert a character to its numeric version (ASCII value/Unicode **codepoint**) with `char->integer`. Reverse: `integer->char`.

```
> (char->integer #\A)
```

```
65
```

```
> (integer->char 65)
```

```
#\A
```

Convert between strings and lists of characters with `string->list` and `list->string`.

```
> (string->list "Howdy thar!")
```

```
(#\H #\o #\w #\d #\y #\space #\t #\h #\a #\r #\!)
```

```
> (list->string '(\a \b \c))
```

```
"abc"
```

Scheme: Strings & I/O

Comparisons [1/5]

We have seen the Scheme numeric comparison operators: = < <= > >=. These can only be used with numbers.

Some Scheme types have their own comparison functions.

```
> (string=? "abc" "def")
```

```
#f
```

```
> (string=? "42" 42)
```

```
ERROR
```

```
> (string<? "abc" "def")
```

```
#t
```

Also: string<=? string>? string>=?
char=? char<? char<=? char>? char>=?

Scheme: Strings & I/O Comparisons [2/5]

There are several kinds of equality in Scheme.

The simplest is `eq?`, which means “same location in memory”.

```
> (eq? '() '())
```

```
#t
```

```
> (eq? 2 2)
```

IMPLEMENTATION-DEPENDENT

```
> (define a '(1 2))
```

```
> (eq? a '(1 2))
```

```
#f
```

```
> (define b a)
```

```
> (eq? a b)
```

```
#t
```

Scheme: Strings & I/O Comparisons [3/5]

Next is `eqv?`, which means “same primitive value”.

```
> (eqv? 2 2)
```

```
#t
```

```
> (eqv? 2 2.0)
```

```
#f
```

```
> (define a '(1 2))
```

```
> (eqv? a '(1 2))
```

```
#f
```

```
> (eqv? "ab" "ab")
```

IMPLEMENTATION-DEPENDENT

Lists and strings are
not primitive values.

Scheme: Strings & I/O

Comparisons [4/5]

Then there is `equal?`, which does the following:

- If the types are different, then return `#f`.
- For primitive values (everything we have covered except strings and pairs) of the same type, call `eqv?`.
- For strings, call `string=?`.
- For pairs, recursively call `equal?` on the `cars` & `cdrs`.

```
> (define a '(1 2))
```

```
> (equal? a '(1 2))
```

```
#t
```

```
> (equal? "ab" "ab")
```

```
#t
```

Scheme: Strings & I/O

Comparisons [5/5]

`equal?` mostly does what we usually want, with one caveat. Since it always returns `#f` when the types are different, it can give unexpected results with numbers.

```
> (equal? 2 2.0)  
#f
```

I offer the following rule of thumb.

- Use `=` for numeric equality.
- Use `equal?` for most other kinds of equality.
- If you want the code to indicate what type you are comparing, and flag type errors for other types, then use a type-specific equality function (e.g., `string=?`, `char=?`).
- Use `eq?` or `eqv?` only if you are sure you know what you are doing.

Scheme: Strings & I/O

Console Output [1/2]

Print any value with `display`. String conversion is automatic. No trailing newline is printed. Print a newline with `newline`. Both of these return `void`, which does not print in the REPL.

```
> (display "Howdy thar!")
```

```
Howdy thar!
```

```
> (newline)
```

```
> (display #\A)
```

```
A
```

```
> (display '(42 #t (300)))
```

```
(42 #t (300))
```

```
> (display +)
```

```
#<procedure: +>
```

Scheme: Strings & I/O

Console Output [2/2]

To do multiple I/O calls in a single expression, use `begin`. This takes any number of arguments, evaluates them all, in order, and returns the value of the last one.

```
> (begin
    (display "dog")
    (display "food")
    (display "love")
  )
dogfoodlove
```

`begin` takes arbitrary expressions, not just those that do I/O.

Scheme: Strings & I/O

Console Interaction [1/5]

Read a line from the console with `read-line`. This takes no parameters. It returns the typed-in line with no trailing newline.

```
> (begin (display "Type something: ") (read-line))
```

```
Type something: Hello there!
```

```
Hello there!
```

Typed by user

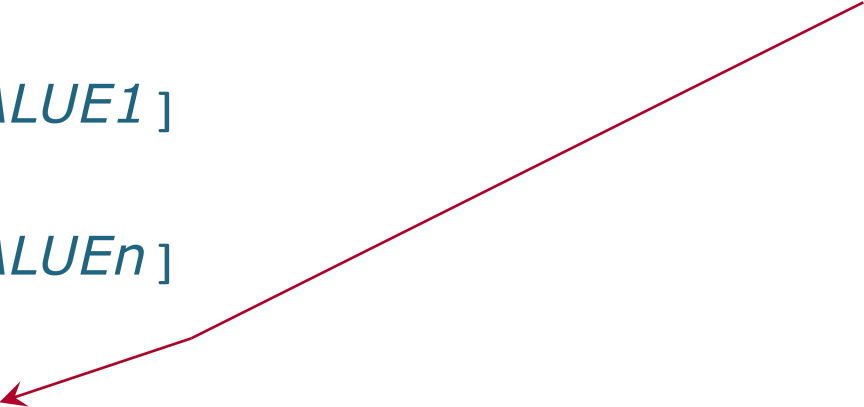
Scheme: Strings & I/O

Console Interaction [2/5]

How can we set a *local* variable to the return value of `read-line` in a procedure?

We use `let`.

```
(let ; Locally bind vars to values in the expression
  (
    [VARIABLE1 VALUE1]
    ...
    [VARIABLEn VALUEn]
  )
  ( EXPRESSION )
)
```



Scheme: Strings & I/O

Console Interaction [3/5]

TO DO

- Write a procedure to prompt for input, read a line, and then print it, with some explanation ("Here is what you typed:").

See string.scm.

Scheme: Strings & I/O

Console Interaction [4/5]

We can repeat using a recursive procedure.
Alternatively, use `let` with a *name*.

```
(let NAME
  (
    [ VARIABLE1 VALUE1 ]
    ...
    [ VARIABLEn VALUEn ]
  )
  ( EXPRESSION )
)
```

Within *EXPRESSION*, we can use *NAME* as a procedure taking *n* arguments. These become the values of *VARIABLE1* through *VARIABLEn* in that invocation of the procedure.

Scheme: Strings & I/O

Console Interaction [5/5]

TO DO

- Write a procedure that reads a series of numbers, until a blank line is entered, printing a running total after each. It should also respond in a reasonable way if the user types something other than a number.

See `string.scm`.